# Chordal-GCN: Exploiting sparsity in training large-scale graph convolutional networks

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#### Abstract

Despite the impressive success of graph convolutional networks (GCNs) on numerous applications, training on large-scale sparse networks remains challenging. Current algorithms require large memory space for storing GCN outputs as well as all the intermediate embeddings. Besides, most of these algorithms involves either random sampling or an approximation of the adjacency matrix, which might unfortunately lose important structure information. In this paper, we propose Chordal-GCN for semi-supervised node classification. The proposed model utilizes the *exact* graph structure (*i.e.*, without sampling or approximation), while requires limited memory resources compared with the original GCN. Moreover, it leverages the sparsity pattern as well as the clustering structure of the graph. The proposed model first decomposes a large-scale sparse network into several small dense subgraphs (called *cliques*), and constructs a clique tree. By traversing the tree, GCN training is performed clique by clique, and connections between cliques are exploited via the tree hierarchy. Furthermore, we implement Chordal-GCN on large-scale datasets and demonstrate superior performance.

# 1 Introduction

Graph convolutional network (GCN) [KW17] is a generalization of convolutional neural networks (CNNs) [LB98] to the graph structure. For a given node, the graph convolution operation aggregates the embeddings (features) of its neighbors, followed by a non-linear transformation. By stacking multiple graph convolutional layers, one can learn node representations by utilizing features of its distant neighborhood. The original GCN model, as well as its numerous variations, has shown great success in a variety of applications, including semi-supervised node classification [KW17], inductive node embedding [HYL17], link prediction [vKW17], and knowledge graphs [SKB<sup>+</sup>18].

Despite the success of GCNs, training GCNs on large-scale graphs remains challenging due to the memory issue: we need to store all the parameters and outputs of GCN. Thus, the memory space scales linearly in the size of graph while quadratically in the feature dimension [CLS<sup>+</sup>19, ZHW<sup>+</sup>19]. This prevents applications of GCN on many real-world networks, including social networks, recommender systems, and knowledge graphs, where the graphs usually contain millions or even billions of nodes.

Methods aimed at large-scale training have been proposed and can be roughly divided into two categories: (1) sampling-based methods and (2) clustering-based methods. For sampling-based methods, only a few neighbors for every node will be sampled in every GCN layer, and thus the size of intermediate embeddings for every layer will be reduced for each mini-batch. Works in this track include [HYL17, CMX18, CZS18, ZHW<sup>+</sup>19]. However, ignorance of some neighbors might lose important structure information, which is the main drawback of all the sampling methods. Another

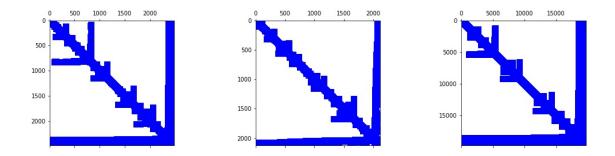


Figure 1: Adjacency matrices for the largest connected component (LCC) of cora (left), citeseer (middle), pubmed (right). The LCCs contain the majority of the graph, so deleting other nodes do not affect the overall arrow pattern. Similar patterns exist in many real-world networks, as also observed in [DBB19].

direction of research notices the sparsity of real-world networks and exploits the clustering structure of the graph. For example, Cluster-GCN [CLS<sup>+</sup>19] separates the graph into several clusters, and in every iteration of training, only one or a few clusters are picked to calculate the stochastic gradient for the mini-batch. However, Cluster-GCN ignores all the inter-cluster links, which are not negligible in many real-world networks. For example, Figure 1 shows the sparsity pattern of three citation networks. We first rearrange the vertices via an approximate minimum degree (AMD) ordering algorithm [ADD96], and then observe a nice arrow pattern in the adjacency matrices. This indicates the existence of some *highly-cited papers* (in the right bottom corner), which have impacts on multiple communities. In Cluster-GCN, these highly-cited papers are randomly put into one community and the adjacency matrix is approximated by a block-diagonal matrix; *i.e.*, Cluster-GCN removes the fletching part of the arrow patterns in Figure 1. This approximation ruins the beautiful arrow pattern, and thus ignores the multi-community influence of some seminal papers.

The above difficulties can be easily tackled by the *Chordal-GCN*, a novel clustering-based method for the semi-supervised node classification task. Recall that Cluster-GCN ignores all the inter-cluster links and trains each cluster separately; comparatively, in Chordal-GCN, we keep all the links, train every cluster separately, and at the same time capture the connections between clusters by an additional loss term. This *partially separable* training process can be achieved with the help of chordal sparsity theory [VA15]: we first decompose a large-scale sparse graph into several small dense subgraphs, and we construct a tree of which the nodes are the subgraphs. Note that two subgraphs are adjacent in the tree if they share some vertices. In the training process, we traverse the tree from leaf to root; and when training on a certain subgraph, we minimize the usual GCN loss, plus an additional term called *consistency loss*. With the consistency loss, messages in the children subgraphs can be passed to their parent, and thus the relationship between subgraphs is leveraged via the hierarchy of the tree. Therefore, Chordal-GCN exploits the sparsity pattern as well as the clustering structure of the graph without any approximation or random sampling, while requires similar memory space to Cluster-GCN.

Our contribution is summarized as follows:

• We propose Chordal-GCN for semi-supervised node classification on large-scale sparse networks. The proposed model fully exploits the *exact* graph structure, while requires limited memory usage on large-scale graphs (much smaller than the original GCN).

- Chordal-GCN is able to train a large-scale graph in a partially separable manner; *i.e.*, in every iteration, the training is performed on a subgraph, and the connections between subgraphs are handled by a consistency loss.
- We analyze the memory and time complexity of Chordal-GCN and compare them with other state-of-the-art GCN models. Also, we evaluate the performance of Chordal-GCN on benchmark datasets and demonstrate superior performance in large-scale datasets.

# 2 Background

In this section, we review the basic concepts of graph convolutional networks (GCNs), and introduce the definition and some fundamental properties of chordal graphs.

#### 2.1 Semi-supervised node classification with GCNs

Suppose G = (V, E) is an undirected graph with |V| = n and the  $n \times n$  symmetric matrix A is the adjacency matrix. Every node k is associated with a feature vector  $x_k \in \mathbb{R}^d$ , and all the feature vectors are stored in the rows of  $X \in \mathbb{R}^{n \times d}$ . An *L*-layer graph convolutional network (GCN) [KW17] has the layer-wise propagation rule:

$$H^{(l+1)} = \sigma_l(\bar{A}H^{(l)}W^{(l)}),$$

where  $W^{(l)} \in \mathbb{R}^{d_l \times d_{l+1}}$  is a weight matrix. The matrix  $\bar{A} = \tilde{D}^{-1/2} \tilde{A} \tilde{D}^{-1/2}$  is the normalized adjacency matrix, where  $\tilde{A} = I + A$  and  $\tilde{D} \in \mathbb{R}^{n \times d}$  is diagonal with  $\tilde{D}_{ii} = \sum_j \tilde{A}_{ij}$ . The matrix  $H^{(l)} \in \mathbb{R}^{n \times d_l}$  are the activations in the *l*-th layer  $(H^{(0)} = X \text{ and } d_0 = d)$ . The function  $\sigma_l$  is the activation function for layer *l*. The output of an *L*-layer GCN is an  $n \times p$  matrix  $H^{(L)}$ , of which the *k*-th row is the predicted label for node *k*. Later in the paper, we will train GCN using only a subgraph  $G(\gamma)$ , so the notation  $y_k^{(\text{pred})}(W; X_{\gamma,:}, \bar{A}_{\gamma\gamma})$  is used to indicate the predicted label of node *k*, using the coefficient matrices  $W = (W^{(0)}, \ldots, W^{(L-1)})$ , the submatrix  $\bar{A}_{\gamma\gamma}$ , and the corresponding feature vectors  $X_{\gamma,:}$ .

Semi-supervised node classification is a popular application of GCN. When applying GCN for this task, we minimize the loss function

$$\mathcal{L}^{\langle \text{gcn} \rangle}(W; X, \bar{A}) = \sum_{k \in V^1} \ell_1 \big( y_k^{\langle \text{true} \rangle}, y_k^{\langle \text{pred} \rangle}(W; X, \bar{A}) \big), \tag{1}$$

to learn the weight matrices  $W = (W^{(0)}, \ldots, W^{(L-1)})$ . In the formula,  $y_k^{\text{(true)}}$  is the given true label for node k, and  $V^l \subseteq V$  is the subset of vertices with given true labels. The function  $\ell_1(\cdot, \cdot)$  is a loss function, and usually it is the cross entropy loss.

#### 2.2 Preliminaries on chordal graphs

Chordal sparsity has been a classical topic in graph theory and found useful applications in various fields, including database theory [BFMY83], probabilistic networks [Pea88, CDLS99, Dar09], linear algebra [Ros70], combinatorial optimization [Gol04, Gav72], semidefinite optimization [ADV10, ADV13]. In this paper, we will use chordal sparsity theory to exploit the clustering structure. So now we show that a chordal graph can be separated into cliques, and that these cliques form a clique tree with desirable properties.

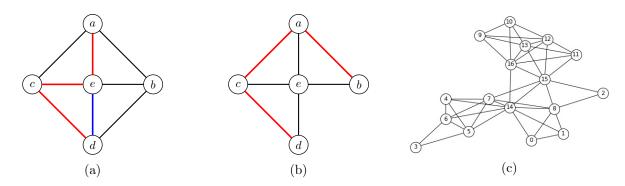


Figure 2: (a) An example of non-chordal graph. The edge  $\{d, e\}$  is a chord in the path (a, e, c, d). This graph is non-chordal because the cycle (a, b, d, c, a) is chordless. (b) An example of chordal graph. The path (b, a, c, d) is chordless. (c) A chordal graph with 17 vertices.

**Basic graph theory.** An undirected graph is a pair G = (V, E), with |V| = n and E a subset of  $\{\{v, w\} \mid v, w \in V\}$ . The elements of V are the vertices of the graph, the elements of E its edges. Two vertices v and w are adjacent if  $\{v, w\} \in E$ . The neighborhood or adjacency set  $\operatorname{adj}(v)$  of a vertex v is the set of vertices adjacent to it:  $\operatorname{adj}(v) = \{w \mid \{v, w\} \in E\}$ , and the degree of a vertex v is the number of its neighbors:  $\operatorname{deg}(v) = |\operatorname{adj}(v)|$ . The adjacency matrix A of the graph G is a  $|V| \times |V|$  symmetric matrix with elements  $A_{ij} = 1$  if  $\{v_i, v_j\} \in E$  and zero otherwise.

**Ordered undirected graphs.** An ordering of an undirected graph G = (V, E) is a numbering of its vertices. We may also use the notation  $\sigma: \{1, 2, ..., n\} \to V$  to denote a sequence of vertices  $\sigma = (\sigma(1), \ldots, \sigma(n))$ . We refer to  $\sigma^{-1}(v)$  as the *index* of vertex v. An ordered graph, denoted as  $G_{\sigma} = (V, E, \sigma)$ , is an undirected graph G = (V, E) with an ordering  $\sigma$ . Thus, the higher neighborhood and lower neighborhood can be defined as

$$adj^+ = \{ w \in adj(v) \mid \sigma^{-1}(w) \ge \sigma^{-1}(v) \}, \quad adj^- = \{ w \in adj(v) \mid \sigma^{-1}(w) \le \sigma^{-1}(v) \}.$$

It is easy to identify the higher and lower adjacency set in the matrix representation form of an ordered graph. The indices of the elements of  $\operatorname{adj}^{-}(v)$  are the column indices of the entries in row  $\sigma^{-1}(v)$  of the matrix; the indices of the elements of  $\operatorname{adj}^{+}(v)$  are the row indices of the entries in column  $\sigma^{-1}(v)$ . We also use the notation  $\operatorname{col}(v)$  and  $\operatorname{row}(v)$  for the closed higher (and lower) neighborhoods

$$\operatorname{col}(v) = \{v\} \cup \operatorname{adj}^+(v), \qquad \operatorname{row}(v) = \{v\} \cup \operatorname{adj}^-(v).$$

**Cliques.** A complete graph is a graph in which every pair of distinct vertices is adjacent. A clique  $\gamma \subseteq V$  of a graph G is a subset of vertices  $\gamma \subseteq V$  that induces a maximal complete subgraph. The graph induced by  $\gamma$  is denoted  $G(\gamma) = (W, E(\gamma))$  and has edge set  $E(\gamma) = \{\{v, w\} \in E \mid v, w \in \gamma\}$ . One does not distinguish between a subset of the vertices and the subgraph induced by it.

**Chordal graph and perfect elimination ordering.** A *path* between  $v_0$  and  $v_k$  is a sequence of distinct vertices  $(v_0, v_1, v_2, \ldots, v_k)$  with  $\{v_i, v_{i+1}\} \in E$  for  $i = 0, 1, \ldots, k - 1$ . For example, the graph in Figure 2a has a path (a, e, c, d). A *chord* in a path of an undirected graph is an edge between non-consecutive vertices on the path. For example, in Figure 2a, the edge  $\{d, e\}$  is a chord

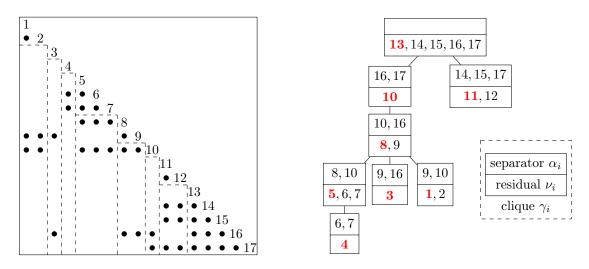


Figure 3: Left. The (symmetric) adjacency matrix of the graph G shown in Figure 2c. (For simplicity we only show the lower triangular part.) A bullet in the (i, j)th entry means nodes i and j are adjacent. The dashed lines separate the supernodes (clique residuals) in the clique tree. Right. The corresponding postordered clique tree  $T^c$ . Every node in  $T^c$  is a clique in G, and the red index i in every clique indicates the clique representative. The dashed box shows an example clique.

in the path (a, e, c, d), but in Figure 2b, the path (b, a, c, d) is chordless. An undirected graph is *chordal* if every cycle of length greater than three has a chord. Figure 2b and 2c show two examples of chordal graphs, while the graph in Figure 2a is non-choral. For a chordal graph G = (V, E), we can always find an ordering  $\sigma$  such that the higher neighborhood of every vertex is complete, *i.e.*,

$$v, w \in \operatorname{adj}^+(u) \implies \{v, w\} \in E.$$

This means that the vertex set col(v) induces a clique for any vertex  $v \in V$ . This ordering is called the *perfect elimination ordering*, and a graph is chordal if and only if it has a perfect elimination ordering [FG65].

Supernodal elimination tree. Given a chordal graph  $G_{\sigma} = (V, E, \sigma)$  with perfect elimination ordering  $\sigma$ , we can always construct a supernodal elimination tree  $T^c$ . It is a rooted tree, of which the vertices are cliques of the sparsity graph  $G_{\sigma}$ . The root of  $T^c$  is the clique that contains the the node  $\sigma(n) \in V$ . Every non-rooted clique  $W = \operatorname{col}(v)$  can be partitioned into clique separators and clique residuals. The clique separator is the intersection of the clique with its parent clique in  $T^c$ , and the clique residual is also called the (maximal) supernode, which contains the node  $v \in V$  and is denoted by  $\operatorname{snd}(v)$ . All the supernodes form a partition of V. Since by our assumption, the node v in the clique  $W = \operatorname{col}(v)$  has the least order, we call v the representative of the supernode and the set of all representatives  $V^c$ .

**Postordered supernodal elimination tree.** For a given supernode partition  $\{\operatorname{snd}(v) \mid v \in V^c\}$ , we can always find a *topological postordering*  $\sigma$ , *i.e.*, an ordering with the following two properties.

• The elements of each supernode  $\operatorname{snd}(v)$  are numbered consecutively: if  $\sigma^{-1}(v) = i$  and  $n_v = |\operatorname{snd}(v) - 1|$ , then  $\operatorname{snd}(v) = \{\sigma(i), \sigma(i+1), \dots, \sigma(i+n_v)\}$ .

• The ordering  $\sigma$  defines a topological ordering of the representative vertices in the supernodal elimination tree  $T^c: \sigma^{-1}(v) \leq \sigma^{-1}(p(v))$  if  $v \in V^c$  and p(v) is the representative of the parent of  $\operatorname{snd}(v)$  in  $T^c$ .

We refer to a supernode elimination tree with certain topological postordering  $\sigma$  as postordered supernodal elimination tree. Given a postordered supernodal elimination tree with a topological postordering  $\sigma$ , we define the following index sets:

- For i = 1, ..., n, define the index set  $\gamma_i = col(i)$ , *i.e.*, the index set  $\gamma_i$  contains the row indices of lower-triangular nonzero elements in column *i*.
- For  $i \in V^c$ , define the index set  $\nu_i = \operatorname{snd}(i)$ , *i.e.*,  $\nu_i = \{i, i+1, \ldots, i+n_i\}$  if  $n_i = |\operatorname{snd}(i)| 1$ .
- For  $i \in V^c$ , define the index set  $\alpha_i = \operatorname{col}(i) \backslash \operatorname{snd}(i) = \gamma_i \backslash \nu_i$ . The index set  $\alpha_i$  is empty if *i* is the root of the supernodal elimination tree.

In short, if  $\sigma(i) \in V^c$  is a representative, then we denote the clique  $\operatorname{col}(i) = \gamma_i = \nu_i \cup \alpha_i$  where  $\nu_i = \operatorname{snd}(i)$  is the supernode and  $\alpha_i$  is the clique separator. Figure 3 shows a supernode partition of a chordal sparsity pattern with a topological postordering. (This example is adapted from [VA15].) In this example,

$$\gamma_5 = \{5, 6, 7, 8, 10\}, \quad \nu_5 = \{5, 6, 7\}, \quad \alpha_5 = \{8, 10\}$$

In the end we emphasize that knowing the above properties of the supernodal elimination tree is not enough to construct one from a given chordal sparsity graph. However, there exist sophisticated but efficient (in O(|E|) algorithms to generate a postordered supernodal elimination tree  $T^c$ , together with the parent function p(v) and all the index sets (*i.e.*,  $\gamma_i$ ,  $\nu_i$ , and  $\alpha_i$  for all  $i \in V^c$ ) [LPP89, PS90]. This process is called *chordal decomposition* or *clique decomposition*; see [VA15] for a more detailed survey on chordal graphs.

# 3 Chordal-GCN

Now we are ready to utilize the chordal decomposition technique in training large-scale sparse networks. We note that most real-world networks are indeed sparse, but not chordal. So we first introduce some data preprocessing steps, followed by the detailed description of Chordal-GCN. We also analyze the time and memory complexity of our model, and compare it with other GCN models.

#### 3.1 Preprocessing

**Reordering and chordal extension.** Chordal extension describes the idea of transforming a non-chordal graph G = (V, E) into a chordal graph G' = (V, E') by adding edges, *i.e.*,  $E \subseteq E'$ . Unfortunately, to find the minimum fill-in is claimed to be an NP-complete problem [RTL76, Yan81], and is not always desirable due to algorithm complexity. However, given an ordering of vertices, it is practical to find a *minimal* chordal extension [Oht76, Ros74]. (It is called minimal if the removal of any added edge results in a non-chordal graph.) That is to say, every minimal chordal extension is associated with certain vertex ordering. Thus, the reordering of vertices is important, especially when we prefer a sparse extended graph. So in this work, we use an approximate minimum degree (AMD) ordering algorithm [ADD96] to first *reorder* the vertices, and then find a minimal chordal extension.

Algorithm 1 Chordal-GCN.

- 1: Input. A normalized adjacency matrix  $\overline{A}$ , a postordered supernodal elimination tree  $T^c$ , feature matrix X, and true labels  $y_k^{\text{(pred)}}$  for  $k \in V^1$ . 2: **Output.** The coefficient matrices  $W = (W^{(0)}, \dots, W^{(L-1)})$ .
- 3: Randomly initialize  $W^{[0]}$ .
- for t = 0, 1, 2, ... do 4:
- for  $i \in V^c$  in topological order **do** 5:
- Compute the loss by (4), *i.e.*, 6:

$$\mathcal{L}_{i}(W) = \sum_{k \in V^{1} \cap \gamma_{i}} \ell_{1} \left( y_{k}, y_{k}^{\langle \text{pred} \rangle}(W; X_{\gamma_{i},:}, \bar{A}_{\gamma_{i}\gamma_{i}}) \right) + \sum_{j \in \text{ch}(i)} \sum_{k \in \alpha_{j}} \ell_{2} \left( y_{k}^{\langle \text{pred} \rangle}(W; X_{\gamma_{i},:}, \bar{A}_{\gamma_{i}\gamma_{i}}), y_{k}^{\langle \text{pred} \rangle}(W^{[t]}; X_{\gamma_{j},:}, \bar{A}_{\gamma_{j}\gamma_{j}}) \right).$$

$$(2)$$

Perform a gradient descent step  $W^{[t+1]} = W^{[t]} - \eta \nabla \mathcal{L}_i(W^{[t]}).$ 7: end for 8:

9: end for

Clique merging. Often, the resulting extension contains many small cliques and it would be more efficient to merge some neighboring cliques, as reported in [AG89]. In this work, traversing the tree from leaf to root, we greedily merge clique  $\gamma_i$  with its parent  $\gamma_{p(i)}$  if  $|\nu_i| \leq \tau$  and  $|\nu_{p(i)}| \leq \tau$ where  $\tau$  is a pre-defined threshold. We denote the resulting graph as G'' = (V, E'').

**Chordal decomposition.** Given the graph G'', we can now construct the postordered supernodal elimination tree  $T^{c}$  (together with all the index sets  $(\gamma_{i}, \nu_{i}, \text{ and } \alpha_{i})$ , and the parent function p(i)). With a little abuse of notation, we also use  $T^{\rm c}$  to indicate the postordered supernodal elimination tree of G. In other words, the preprocessing steps reorder and partition the vertex set V, which helps us identify the clique  $\gamma_i$  and the corresponding induced subgraph  $G(\gamma_i)$ . Although chordal extension and clique merging add edges into the graph, it will not affect the training process of GCN because we will use the subgraph  $G(\gamma_i)$ , not  $G''(\gamma_i)$ .

Lastly, we emphasize that, although  $G(\gamma_i)$  no longer represents a complete subgraph, we still refer to  $\gamma_i$  as *clique* to distinguish from *clusters* or *communities*: in the chordal decomposition, one node  $k \in V$  is allowed to exist in multiple cliques while in community detection or graph clustering, one node appears in only one cluster.

#### 3.2Chordal-GCN: meta algorithm

With the postordered supernodal elimination tree  $T^{\rm c}$  in hand, we are able to train a large-scale sparse graph in a *partially separable* manner, *i.e.*, clique by clique. Suppose we traverse the tree  $T^{\rm c}$  in the topological order, *i.e.*, from leaf to root. At clique  $\gamma_i$ , the loss function consists of two parts. The first part is the usual GCN loss defined in (1), but using only the induced subgraph  $G(\gamma_i)$  instead of the entire G:

$$\mathcal{L}^{\langle \text{gcn} \rangle}(\cdot; X_{\gamma_i,:}, \bar{A}_{\gamma_i \gamma_i}) = \sum_{k \in V^1 \cap \gamma_i} \ell_1 \big( y_k^{\langle \text{true} \rangle}, y_k^{\langle \text{pred} \rangle}(\cdot; X_{\gamma_i,:}, \bar{A}_{\gamma_i \gamma_i}) \big),$$

where  $X_{\gamma_i,:}$  takes the rows of X indexed by  $\gamma_i$  and  $\bar{A}_{\gamma_i\gamma_i}$  is the submatrix of  $\bar{A}$  with rows and columns indexed by  $\gamma_i$ .

The second part of the loss is the consistency loss. When we arrive at clique  $\gamma_i$ , we have traversed all its children in the tree  $T^c$ . This means that some nodes in  $\gamma_i$  have already been visited; namely, those clique separators of  $\gamma_i$ 's children in  $T^c$ . For clarity, we denote by ch(i) the set of children representatives of  $\gamma_i$ ; *i.e.*,  $j \in ch(i)$  if j is a representative (*i.e.*,  $j \in V^c$ ) and  $\gamma_j$  is a child of  $\gamma_i$  in  $T^c$ . Then all the nodes in  $\alpha_j$  for  $j \in ch(i)$  have already been visited when we arrive at  $\gamma_i$ . Thereby, for  $k \in \alpha_j$  and  $j \in ch(i)$ , we hope that the predicted label  $y_k^{\langle \text{pred} \rangle}(\cdot; X_{\gamma_i,:}, \bar{A}_{\gamma_i \gamma_i})$  using the current clique  $\gamma_i$  is consistent with the previous prediction  $y_k^{\langle \text{pred} \rangle}(\tilde{W}; X_{\gamma_j}, \bar{A}_{\gamma_j \gamma_j})$  using clique  $\gamma_j$  and the most recent parameter  $\tilde{W}$ . Hence, the consistency loss for training clique  $\gamma_i$  can be formulated as

$$\mathcal{L}^{\langle \text{cons} \rangle}(\cdot; X_{\gamma_i,:}, \bar{A}_{\gamma_i \gamma_i}) = \sum_{j \in \text{ch}(i)} \sum_{k \in \alpha_j} \ell_2 \big( y_k^{\langle \text{pred} \rangle}(\cdot; X_{\gamma_i,:}, \bar{A}_{\gamma_i \gamma_i}), y_k^{\langle \text{pred} \rangle}(\tilde{W}; X_{\gamma_j,:}, \bar{A}_{\gamma_j \gamma_j}) \big), \tag{3}$$

where  $\ell_2(\cdot, \cdot)$  is a loss function, for example, the KL divergence. Note that  $\tilde{W}$  is not treated as the variable of  $\mathcal{L}^{\langle \text{cons} \rangle}$ ; *i.e.*, the term  $y_k^{\langle \text{pred} \rangle}(\tilde{W}; X_{\gamma_j,:}, \bar{A}_{\gamma_j \gamma_j})$  is considered constant when we take the gradient of  $\mathcal{L}^{\langle \text{cons} \rangle}$ .

Therefore, when training the clique  $\gamma_i$ , we compute the total loss

$$\mathcal{L}_{i}(\cdot) = \mathcal{L}^{\langle \text{gcn} \rangle}(\cdot; X_{\gamma_{i},:}, \bar{A}_{\gamma_{i}\gamma_{i}}) + \mathcal{L}^{\langle \text{cons} \rangle}(\cdot; X_{\gamma_{i},:}, \bar{A}_{\gamma_{i}\gamma_{i}}), \tag{4}$$

and perform a gradient descent step  $W^+ = W - \eta \nabla \mathcal{L}_i(W)$  where  $\eta$  is the learning rate. The whole algorithm is summarized in Algorithm 1. Note that the gradient descent step in line 7 of Algorithm 1 can be replaced with any other accelerated version, for example, Adam [KB15].

#### 3.2.1 Connection to the original GCN

We note that Algorithm 1 is essentially the mini-batch gradient descent method applied to the loss

$$\mathcal{L}^{\langle \operatorname{cgcn} \rangle}(W) = \sum_{i \in V^{c}} \sum_{k \in V^{1} \cap \gamma_{i}} \ell_{1} \left( y_{k}^{\langle \operatorname{true} \rangle}, y_{k}^{\langle \operatorname{pred} \rangle}(W; X_{\gamma_{i},:}, \bar{A}_{\gamma_{i}\gamma_{i}}) \right) + \sum_{i \in V^{c}} \sum_{j \in \operatorname{ch}(i)} \sum_{k \in \alpha_{j}} \ell_{2} \left( y_{k}^{\langle \operatorname{pred} \rangle}(W; X_{\gamma_{i},:}, \bar{A}_{\gamma_{i}\gamma_{i}}), y_{k}^{\langle \operatorname{pred} \rangle}(W; X_{\gamma_{j},:}, \bar{A}_{\gamma_{j}\gamma_{j}}) \right),$$

$$(5)$$

with  $\gamma_i$  taken as the mini-batch. Clearly, it is equal to the original GCN loss (1) if we have only one clique (*i.e.*,  $V^c = \{1\}$  and thus  $\gamma_1 = V$ ). Also, when  $\bar{A}$  is block diagonal, Chordal-GCN is equivalent to Cluster-GCN because the consistent loss (3) vanishes.

However, in the general case, without any assumption on the graph structure or the GCN structure, it is impossible to show the equivalence between Chordal-GCN and GCN. From another view of point, the non-convex functions (1) and (5) don't share any global minimizer. For example, consider the chordal graph with two overlapping dense principal submatrices. We can easily choose the set of true labels  $y^{\text{(true)}}$  such that (1) and (5) don't have the same minimizer.

#### 3.2.2 Time and memory efficiency

Following [CLS<sup>+</sup>19], we report the time and memory complexity in Table 1. Although our method has the same time complexity as Cluster-GCN in big-O notation, Chordal-GCN needs an extra

Complexity	GCN [KW17]	GraphSAGE [HYL17]	FastGCN [CMX18]	VR-GCN [CZS18]	Cluster-GCN [CLS <sup>+</sup> 19]	Chordal-GCN
Time Memory	$\begin{array}{c} O(L\ A\ _0d+Lnd^2) \\ O(Lnd+Ld^2) \end{array}$	$\begin{array}{c} O(r^L n d^2) \\ O(br^L d + L d^2) \end{array}$	$O(Lrnd^2)$ $O(Lbrd + Ld^2)$	$O(L  A  _0d + (L+r^L)nd^2)$ $O(Lnd + Ld^2)$	$O(L  A  _0d + Lnd^2)$ $O(Lc_1d + Ld^2)$	$O(L  A  _0d + Lnd^2)$ $O(Lc_2d + Ld^2)$

Table 1: Time and memory complexity of GCN training algorithms. L is the number of layers, n is the number of nodes,  $||A||_0$  is the number of nonzeros in the adjacency matrix A, and d is the number of features. For simplicity we assume the number of features is fixed for all layers, *i.e.*,  $d_l = d$  for l = 1, 2, ..., L. For SGD-based methods, b is the batch size and r is the number of sampled neighbors per node. Also,  $c_1$  is the maximum cluster size in Cluster-GCN while  $c_2$  is the maximum clique size in our method. For memory complexity,  $Ld^2$  is used to store the parameters  $\{W^{(l)}\}_{l=0}^{L-1}$  and the other term is for storing embeddings  $\{H^{(l)}\}_{l=1}^{L}$ . For simplicity we omit the memory for storing the graph or subgraphs since they are fixed and usually not the main bottleneck.

forward propagation on the children cliques in order to compute the predicted labels

$$y_k^{\text{(pred)}}(W^{[t]}; X_{\gamma_j,:}, \bar{A}_{\gamma_j\gamma_j})$$

in (2). The extra forward propagation step would be the only price paid for single node being in multiple subgraphs. As we will see in the numerical experiments, this additional cost is minimal and worthwhile, especially in citation networks with an arrow pattern.

In addition, the memory bottleneck for Chordal-GCN and Cluster-GCN depend on the maximum cluster size and maximum clique size, respectively. Since these two numbers depend totally on the hyperparameters, *i.e.*, number of clusters in METIS algorithm and the merging threshold  $\tau$ , it is difficult to compare analytically. But both are much smaller than the memory usage of the original GCN, which is desirable.

# 4 Experiments

#### 4.1 Experiment settings

We evaluate our proposed model for the semi-supervised node classification task on four public datasets: cora, citeseer, pubmed, and reddit. We compare Chordal-GCN with the following state-of-the-art GCN training algorithms for comparison: GCN [KW17], GraphSAGE [HYL17], FastGCN (both uniform and importance sampling) [CMX18], VR-GCN [CZS18], and Cluster-GCN [CLS<sup>+</sup>19]. For all the methods we use the same GCN structure (*i.e.*, L = 2), the same label rate, and Adam optimizer [KB15] with learning rate 0.001, dropout rate as 20%, weight decay as zero. For the baselines, we use the implementations provided by the authors, and follow the default parameter settings in these models. Other implementation details are included in appendix A.

For all the methods and datasets, we conduct training for 10 times and take the mean of the evaluation results. Also, we stop training when the validation accuracy does not increase a threshold (0.01) for 10 epoches, and choose the model with the highest validation accuracy as convergent point. We use the following metrics to evaluate the performance of all methods.

- Accuracy: The micro F1 score of the test data at the convergent point.
- Memory usage: The maximum GPU memory occupied by tensors.

• Epoch time and number of epoches: Time to run an epoch and total number of epoches before convergence. Note that in Chordal-GCN, an epoch is a traverse of the clique tree, rather than one clique.

### 4.2 Numerical results

We summarize all the results in Table 2. In the three citation networks (cora, citeseer, and pubmed), Chordal-GCN achieves the highest accuracy among all baseline models. This desirable result is due to the exploitation of the exact sparsity pattern and clustering structure.

In terms of memory space, Chordal-GCN is superior to the original GCN, just as expected. All the sampling-based methods require limited memory space because they reduce the number of intermediate node embeddings. More interestingly, Chordal-GCN uses less memory space than Cluster-GCN (except in reddit). This is because the Cluster-GCN implementation uses a few clusters in one epoch, and thus the memory bottleneck depends on the sum of the selected clusters, rather than the maximum cluster. The abnormal memory cost of Chordal-GCN in reddit dataset is attributed to our clique merging heuristic. The maximal clique in reddit has size 220,069, almost 95% of the vertices. Splitting the graph into more cliques will help reduce the memory requirement, while increases the epoch time: we train smaller but more cliques in one epoch.

Lastly, we emphasize that the comparison of epoch time can never be fair: one epoch in Chordal-GCN means one traverse of the clique tree while in other methods one epoch is one mini-batch. This explains the long epoch time of Chordal-GCN.

# 5 Conclusion

We propose Chordal-GCN for semi-supervised node classification on large-scale sparse networks. The proposed model exploits the sparsity pattern and clustering structure of the graph, and utilizes the *exact* graph structure (*i.e.*, without sampling or approximation). Moreover, the memory usage of Chordal-GCN is limited because the training is performed in a partially separable manner; *i.e.*, clique by clique. Experiment results demonstrate that Chordal-GCN achieves the best test accuracy with much smaller memory cost on benchmark datasets.

# 6 Related work

#### 6.1 Related work on chordal graphs

Chordal graphs have attracted much interest in graph theory and found useful applications in various fields such as database theory [BFMY83], probabilistic networks [Pea88, CDLS99, Dar09], linear algebra [Ros70], combinatorial optimization [Gol04, Gav72], semidefinite optimization [ADV10, ADV13], etc. For a complete survey on chordal graphs, please refer to [VA15].

The perfect elimination ordering of chordal graphs has been studied in 1960s [FG65], and there exist linear-time algorithms to test whether a graph G is chordal by looking for a perfect elimination ordering [RTL76, TY84, Heg06]. The theoretical study of postordered supernodal elimination tree traced back to 1990s [LPP89, PS90], and the notation of the postordered supernodal elimination tree used in the paper is taken from the survey paper [VA15]. A recent python implementation

Dataset	Method	F1 (%)	Mem. (MB)	Epoch time (ms)	# epoches
Cora (2708)	GraphSAGE	78.4	565.12	44.26	11
	FastGCN (uniform)	78.2	28.42	27.10	88
	FastGCN (importance)	83.5	28.42	30.41	84
	VR-GCN	81.6	4.60	14.90	64
	Cluster-GCN	68.2	36.14	62.65	8
	GCN	81.9	20.23	5.35	119
	Chordal-GCN	84.3	5.69	90.90	49
	GraphSAGE	61.8	932.14	37.97	60
	FastGCN (uniform)	70.2	98.93	111.00	53
Citeseer	FastGCN (importance)	72.2	98.93	104.90	61
(3327)	VR-GCN	71.2	7.01	32.00	35
(3321)	Cluster-GCN	62.8	60.35	82.09	5
	GCN	70.5	52.16	6.42	95
	Chordal-GCN	74.1	8.47	38.91	57
	GraphSAGE	76.8	475.84	43.57	47
	FastGCN (uniform)	77.4	107.20	25.41	52
Pubmed	FastGCN (importance)	78.0	107.20	22.00	47
(19717)	VR-GCN	78.4	14.49	25.21	33
(19717)	Cluster-GCN	71.2	25.35	59.84	10
	GCN	78.0	66.03	7.52	109
	Chordal-GCN	80.2	15.37	55.94	83
	GraphSAGE	93.1	1192.00	185.22	62
	FastGCN (uniform)	92.9	1656.14	15.83	64
Deddit	FastGCN (importance)	93.0	1722.05	21.60	31
Reddit (232965)	VR-GCN	93.2	1054.21	5041.32	32
	Cluster-GCN	96.1	231.32	2407.61	43
	GCN	93.9	2992.12	261.81	190
	Chordal-GCN	94.2	2501.23	987.12	410

Table 2: Comparison of Chordal-GCN with original GCN, GraphSAGE, FastGCN with uniform/importance sampling, VR-GCN, and Cluster-GCN. We report the micro F1 score (%), memory usage (MB), time per epoch (ms), and number of epoches. Only GCN and Chordal-GCN utilize the exact graph, while other baseline methods lose some graph structure information. for postordered supernodal elimination tree is due to Andersen and Vandenberghe [AV15], which is used in our implementation.

#### 6.2 Related work on graph convolutional networks

A lot of work has been done to improve the performance of GCN. In the original paper [KW17], the authors compute the batch gradient in the training process, which involves the entire labeled set of nodes  $V^{\rm l}$ . A remedy is to approximate the batch gradient by a stochastic gradient; *i.e.*, in every iteration, one uses a randomly selected induced subgraph instead of the whole graph. However, this gradient is still expensive to compute, due to the large receptive field size. Thus, modification on the normalized adjacency matrix (or Laplacian matrix)  $\overline{A}$  has been made for a more efficient evaluation of the matrix-vector product  $\overline{AH}^{(l)}$ . These works can be roughly divided into two categories: the sampling methods, and the clustering methods.

The sampling method randomly picks some neighbors for each node  $i \in V$  at layer l. This method dramatically decreases the number of multiplications one needs to perform in computing  $\bar{A}H^{(l)}$  because only very few elements  $\bar{A}_{ij}$  are involved in the calculation. The rules for random selection can either be uniform sampling (or neighborhood sampling) [HYL17], or importance sampling (known as FastGCN [CMX18]), *i.e.*, the probability of picking a node  $i \in V$  is proportional to  $\sum_{j=1}^{n} \bar{A}_{ij}$ . However, these heuristics are only proposed as computational tricks, and there is no theoretical guarantees whether these sampling methods are able to recover the performance of the original GCN. Basically, sampling might lose information contained in the adjacency matrix, and thus the network structure. A remedy for this issue would be variance reduction GCN [CZS18], in which the authors find another balance between the precision and computational efficiency; *i.e.*, only the "history" part  $\bar{A}H_{old}^{(l)}$  is computed precisely (using the full adjacency matrix  $\bar{A}$ ), while the "correction" part  $\bar{A}^{(l)}(H_{new}^{(l)} - H_{old}^{(l)})$  is computed using  $\bar{A}^{(l)}$  chosen via uniform sampling.

Another direction of research tries to exploit the clustering characteristic in some social networks. For example, Cluster GCN [CLS<sup>+</sup>19] separates the underlying graph into several clusters, and simply ignores the inter-cluster links. The resulting adjacency matrix has a nice block-diagonal structure, and then in each iteration of the training process, only a few clusters are chosen randomly to calculate the stochastic gradient for the mini-batch. This idea can be viewed as a special case of chordal GCN since block diagonal pattern is chordal. In addition, every node can only appear once in one cluster (or community) as in cluster GCN while it can appear in different cliques in our model. Our model is more reasonable because in most citation networks, for example, a highly-cited paper usually has extensive influence in multiple communities. So in our model, this highly-cited paper appears in multiple cliques and affects research work in different communities.

The clustering idea can be applied in terms of graph Laplacian. Recall that the number of zero eigenvalues in the graph Laplacian indicates the number of connected components in a graph. So it is reasonable to approximate the Laplacian by a low-rank matrix. This idea is used in the LanczosNet [LZUZ19], in which the eigenvalues are computed via the classical Lanczos process. In addition, the Lovász convolutional network [YNY<sup>+</sup>19] replaces  $\bar{A}$  with another matrix K, which has the same sparsity pattern as  $\bar{A}$  (*i.e.*,  $K_{ij} = 0$  if  $\bar{A}_{ij} = 0$ ) but is low-rank and positive semidefinite.

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# A Implementation details

Hardware and software. We use Tesla V100 GPU, Intel Xeon CPU E5-2698 (2.20GHz) with 500GB. We use PyTorch 1.2.0 in our implementation and Tensorflow 1.13 to test the baselines.

**Datasets.** The cora, citeseer, and pubmed datasets are from https://github.com/tkipf/gcn. The reddit data is from https://snap.stanford.edu/graphsage. We report the data statistics in Table 3.

Dataset	Nodes	Edges	Classes	Features	Label rate
Cora	2,708	5,429	7	1,433	140(5.1%)
Citeseer	3,327	4,732	6	3,703	120~(3.6%)
Pubmed	19,717	44,338	3	500	60~(0.3%)
Reddit	232,965	11,606,919	41	602	152,410~(65.4%)

Table 3: Data statistics

**Preprocessing.** In Chordal-GCN, we use the AMD ordering [ADD96] implemented in the python package CVXOPT [ADV15]. The minimal chordal extension and the chordal decomposition are performed using the python package chompack [AV15]. We also use the clique merging heuristic described in section 3.1. The threshold  $\tau$  is different in different datasets, and mainly depends on the size of the graph. We report the parameters in Table 4.

**GCN structure.** We implement our model with PyTorch. We use a two-layer GCN with fixed parameter size 128.

Dataset	Threshold for merging $(\tau)$	Number of cliques	Maximum clique size	Minimum clique size
cora	100	19	446	125
citeseer	200	6	512	227
pubmed	2,000	8	6,105	2,132
reddit	10,000	126	220,069	1,001

Table 4: Parameters in Chordal-GCN

To calculate the meory usage, we use tf.contrib.memory\_BytesInUse() for TensorFlow and torch.cuda.memory\_allocated() for PyTorch.

# **B** Additional experiments

# B.1 Ablation study

We evaluate the contribution of the consistency loss  $\mathcal{L}^{\langle \text{cons} \rangle}$  (3) here. We train Chordal-GCN with and without  $\mathcal{L}^{\langle \text{cons} \rangle}$  on the cora, citeseer, and pubmed datasets, and plot the validation accuracy

versus the number of epoches. Figure 4 shows that the consistency loss  $\mathcal{L}^{\langle \text{cons} \rangle}$  correctly captures the connection between cliques, and that improve the performance of Chordal-GCN. This explains why our model performs consistently better than Cluster-GCN: Cluster-GCN ignores those intercluster links and thus loses the connections between clusters. In the pubmed dataset, Chordal-GCN without consistency loss is even overfitting: the validation accuracy decreases after several epoches.

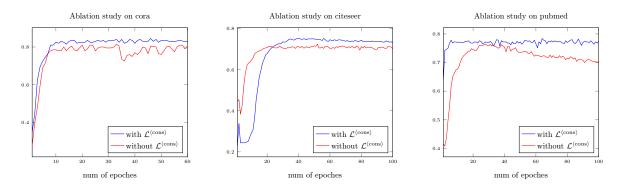


Figure 4: Validation accuracy of Chordal-GCN on the datasets: cora, citeseer, and pubmed. The blue curve shows the validation accuracy with the consistent loss  $\mathcal{L}^{\langle \text{cons} \rangle}$  while the red curve shows that without  $\mathcal{L}^{\langle \text{cons} \rangle}$ .

# B.2 Training deep GCNs

In this section, we study the performance of deeper Chordal-GCN, and explore how time and memory cost scales with GCN depth. As shown in Figure 5, both memory and time cost scale linearly with the number of layers. This result is consistent with theoretical analysis, and shows that Chordal-GCN is applicable to training deeper GCNs.

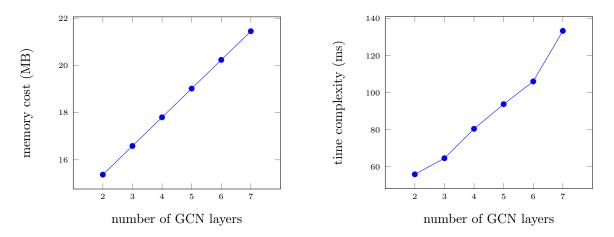


Figure 5: Memory (left) and time (right) versus number of GCN layers. Both curves scale linearly, which is promising and consistent with theoretical analysis.